

# Operational Semantics of $\mu$ Scheme

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April 10, 2024

- Assignment #3 will be out soon

$$\langle e, \xi, \phi, \rho \rangle \Downarrow \langle v, \xi', \phi, \rho' \rangle \quad \langle d, \xi, \phi \rangle \rightarrow \langle \xi', \phi' \rangle$$

$$\langle e, \xi, \phi, \rho \rangle \Downarrow \langle v, \xi', \phi, \rho' \rangle \quad \langle d, \xi, \phi \rangle \rightarrow \langle \xi', \phi' \rangle$$

- $e$ : expression to evaluate
- $\xi$ : global environment
- $\phi$ : function environment (does not change in evaluation!)
- $\rho$ : parameter environment

$$\langle e, \rho, \sigma \rangle \Downarrow \langle v, \sigma' \rangle$$

$$\langle d, \rho, \sigma \rangle \rightarrow \langle \rho', \sigma' \rangle$$

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- $\rho$  maps variables to **locations**  $l$

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- $\rho$  maps variables to **locations**  $l$  (**does not change in evaluation**)
- $l \in \{l_0, l_1, l_2, \dots\}$  represents an **address** in memory



$$\langle e, \rho, \sigma \rangle \Downarrow \langle v, \sigma' \rangle \qquad \langle d, \rho, \sigma \rangle \rightarrow \langle \rho', \sigma' \rangle$$

- $\rho$  maps variables to **locations**  $l$  (**does not change in evaluation**)
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- $\sigma$  is the **memory**: it maps locations to values.

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- $\rho$  maps variables to **locations**  $l$  (does not change in evaluation)
- $l \in \{l_0, l_1, l_2, \dots\}$  represents an **address** in memory
- $\sigma$  is the **memory**: it maps locations to values.
- The memory *can* change (though we haven't used this)

Main challenge: managing state of **closures**

# Variables

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$$\frac{x \in \text{dom}(\rho) \quad \rho(x) \in \text{dom}(\sigma)}{\langle \mathbf{VAR}(x), \rho, \sigma \rangle \Downarrow \langle \sigma(\rho(x)), \sigma \rangle}$$

$$\frac{x \in \text{dom}(\rho) \quad \rho(x) = l \quad \langle e, \rho, \sigma \rangle \Downarrow \langle v, \sigma' \rangle}{\langle \text{SET}(x, e), \rho, \sigma \rangle \Downarrow \langle v, \sigma' \{l \mapsto v\} \rangle}$$

## Local variables

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$\mu$ Scheme allows us to declare local variables

```
(let ((x value-of-x)
      (y value-of-y)
      ...))
  expression-with-x-and-y)
```



Case of one variable declaration:

$$\frac{l_x \notin \text{dom}(\sigma_1) \quad \langle e_x, \rho, \sigma_0 \rangle \Downarrow \langle v_x, \sigma_1 \rangle \quad \langle e, \rho\{x \mapsto l_x\}, \sigma_1\{l_x \mapsto v_x\} \rangle \Downarrow \langle v, \sigma_2 \rangle}{\langle \mathbf{LET}(x, e_x, e), \rho, \sigma_0 \rangle \Downarrow \langle v, \sigma_2 \rangle}$$

# Functions

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$$\langle \text{LAMBDA}(x, e), \rho, \sigma \rangle \Downarrow \langle (\|\text{LAMBDA}(x, e), \rho\|), \sigma \rangle$$

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$$\langle \text{LAMBDA}(x, e), \rho, \sigma \rangle \Downarrow \langle \llbracket \text{LAMBDA}(x, e) \rrbracket, \rho \rrbracket, \sigma \rangle$$

$$\begin{array}{l}
 l_x \notin \text{dom}(\sigma) \quad \langle e, \rho, \sigma \rangle \Downarrow \langle \llbracket \text{LAMBDA}(x, e_c) \rrbracket, \rho_c \rrbracket, \sigma_0 \rangle \\
 \langle e_x, \rho, \sigma_0 \rangle \Downarrow \langle v_x, \sigma_1 \rangle \quad \langle e_c, \rho_c \{x \mapsto l_x\}, \sigma_1 \{l_x \mapsto v_x\} \rangle \Downarrow \langle v, \sigma' \rangle \\
 \hline
 \langle \text{APPLY}(e, e_x), \rho, \sigma \rangle \Downarrow \langle v, \sigma' \rangle
 \end{array}$$

## Supporting multiple arguments/variables

- Check that variables are distinct
- Create multiple fresh locations
- Evaluate arguments/initial values left to right

## Example

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## Evaluate the following expression

```
(let ((f (let ((x 1))  
           (lambda (y) (set x (+ x y))))))  
  (f (f 5)))
```

## First, evaluate the lambda

Let's abbreviate

$$b_f = \text{SET}(x, (\text{APPLY}(\text{VAR}(+), \text{VAR}(x), \text{VAR}(y))))$$

$$e_f = \text{LAMBDA}(y, b_f)$$

$$c_f = \langle e_f, \{x \mapsto l_x\} \rangle$$

Derive

$$D_\lambda = \frac{}{\langle e_f, \{x \mapsto l_x\}, \{l_x \mapsto 1\} \rangle \Downarrow \langle c_f, \{l_x \mapsto 1\} \rangle}$$



Next, evaluate the inner let

$$D_x = \frac{\overline{\langle 1, \emptyset, \emptyset \rangle \Downarrow \langle 1, \emptyset \rangle} \quad l_x \notin \text{dom}(\emptyset) \quad D_\lambda}{\langle \text{LET}(x, 1, e_f), \emptyset, \emptyset \rangle \Downarrow \langle c_f, \{l_x \mapsto 1\} \rangle}$$

## Evaluate the first call

$$\rho_0 = \{f \mapsto c_f\}$$

$$\rho_1 = \{x \mapsto l_x, y \mapsto l_y\}$$

$$\sigma_0 = \{l_f \mapsto c_f, l_x \mapsto 1\}$$

$$\sigma_1 = \{l_f \mapsto c_f, l_x \mapsto 1, l_y \mapsto 5\}$$

$$\sigma_2 = \{l_f \mapsto c_f, l_x \mapsto 6, l_y \mapsto 5\}$$

$$l_y \notin \{l_f, l_x\}$$

...

...

$$D_{f5} = \frac{\frac{\langle \text{VAR}(f), \rho_0, \sigma_0 \rangle \Downarrow \langle c_f, \sigma_0 \rangle}{\dots} \quad \frac{\langle 5, \rho_0, \sigma_0 \rangle \Downarrow \langle 5, \sigma_0 \rangle}{\dots} \quad \frac{\langle b_f, \rho_1, \sigma_1 \rangle \Downarrow \langle 6, \sigma_2 \rangle}{\dots}}{\langle \text{APPLY}(\text{VAR}(f), 5), \rho_0, \sigma_0 \rangle \Downarrow \langle 6, \sigma_2 \rangle}$$

## Evaluate the second call

$$\rho_2 = \{x \mapsto l_x, y \mapsto l'_y\}$$

$$\sigma_3 = \sigma_2\{l'_y \mapsto 6\}$$

$$\sigma_4 = \sigma_3\{l_x \mapsto 12\}$$

...

...

$$D_{f(f5)} = \frac{l'_y \notin \{l_f, l_x, l_y\} \quad \overline{\langle \text{VAR}(f), \rho_0, \sigma_0 \rangle \Downarrow \langle c_f, \sigma_0 \rangle} \quad D_{f5} \quad \overline{\langle b_f, \rho_2, \sigma_3 \rangle \Downarrow \langle 12, \sigma_4 \rangle}}{\langle \text{APPLY}(\text{VAR}(f), \text{APPLY}(\text{VAR}(f), 5)), \rho_0, \sigma_0 \rangle \Downarrow \langle 12, \sigma_4 \rangle}$$

$$D_f = \frac{l_f \notin \emptyset \quad D_x \quad D_{f(f5)}}{\langle \text{LET}(f, \text{LET}(x, 1, e_f), \text{APPLY}(\text{VAR}(f), \text{APPLY}(\text{VAR}(f), 5))), \emptyset, \emptyset \rangle}$$
$$\Downarrow \langle 12, \{l_f \mapsto c_f, l_x \mapsto 12, l_y \mapsto 5, l'_y \mapsto 6\} \rangle$$